

Computer Programming

Timing. Randomization. Hashing

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## Date and time (time.h)

time.h contains structures and functions to measure time

clock\_t and time\_t are real types representing times

**struct** tm holds a broken-down calendar time (sec, min, ... year)

**struct** timespec holds time in seconds and nanoseconds

```
clock_t clock(void);
```

returns (approximation) of processor time used

divide by CLOCKS\_PER\_SEC (usually  $10^6$ ) to get time in seconds

```
int timespec_get(struct timespec *ts, int base);
```

gives time in s and ns since a reference point base (use TIME\_UTC)

```
struct timespec {  
    time_t tv_sec;  
    long tv_nsec;  
};
```

## Measuring time

Place the code to be benchmarked in a loop running many times  
total time: order of seconds (account for limited clock precision)

Ensure compiler doesn't optimize away repetition (check assembly)  
e.g. computing/assigning the same value many times  
may need to use **volatile** specifier for variables  
(forces writing/reading to memory every time, like in source)

Repeat measurements and make an average.

Time may be affected by other running processes, caching, etc.

## Pseudo-random numbers (`stdlib.h`)

Only natural phenomena can be truly random.

Computer uses algorithm to generate numbers  $\Rightarrow$  *pseudo-random*

period of number generator should be high

all bits should appear to be random

Quality of `stdlib` random number generator may not be high

(esp. for lower bits)

Need to use special RNG in cryptography applications.

```
int rand(void);
```

returns an integer in range 0 to `RAND_MAX` (at least  $2^{15} - 1$ )

Re-running program will produce the same sequence of numbers!

$\Rightarrow$  need to initialize state of RNG with a *seed*

```
void srand(unsigned int seed);
```

could use calendar time (seconds) as seed – different in each run

```
e.g. srand((unsigned)time(NULL));
```

# Hashing

*Searching* is a fundamental and widely encountered problem  
find whether an object belongs to a set of objects already stored  
use *maps* (functions) from arbitrary *keys* to *values*  
arrays only work when keys are integers (in a given range)

Idea: find a function  $h$  with an integer value in a restricted range  
(usable as index in an array)

Every object (key)  $x$  is stored in array at index  $h(x)$   
(usually,  $h(x)$  modulo table size)

Objects with different hash value are surely different  
Different objects may have same hash value (collision)

# Hash functions

need to be fast (easily computed), mixing all bytes of the object  
have few *collisions* (esp. for objects with close/related values)

clearly, collisions cannot be avoided if domain larger than range

Examples for strings:

```
for (h=len; len--;) h = ((h<<7) ^ (h<<27)) ^ *s++; // Knuth
for (h=5381; c=*s++; ) h += (h << 5) + c; // Bernstein
for (h=0; c=*s++; ) h = (h<<6) + (h<<16) - h + c; // SDBM
```

Hash functions usually return (32-bit) integers

Cryptographic functions need more stringent properties, and have larger bit width (128-256 bits)

# Open and closed hashing

## *Closed hashing*

if a different object is found at index  $idx=h(x)$ , continue search using a sequence of indices:

sequential:  $idx++$ , linear:  $idx+=i$

with another hash function:  $idx+=h_2(x)$  until element found

when table fills up, objects must be re-hashed

deleted objects must be marked ( $\neq empty$ ) to stop useless search

## *Open hashing*

entry in hash table is (linked) *list* of objects with same hash value

$\Rightarrow$  hashing followed by linear search in (hopefully short) list

need dynamic allocation for list elements

hash table size comparable to element count (avoid long lists)

# Cuckoo hashing (Pagh & Rodler 2001)

constant-time lookup

amortized constant-time insert

Each key may be found in one of *two* locations  
(use two different hash functions)

On collision, displace existing key to 2<sup>nd</sup> location;  
if that location is full, successively displaced

If a cycle is reached, rebuild (larger) table

Works well up to  $\sim 50\%$  fill factor

Arrows in figure show alternate location for a key

