

Computer Programming

Implementing an abstract datatype.  
Linked lists and queues

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## Libraries and abstract datatypes

Use of (standard) library so far:

we know a *function prototype* (declaration), e.g.

```
FILE *fopen(const char *fname, const char *mode);
```

*declaration* is included from *header file* `#include <stdio.h>`

we do *not know or need the source* code for fopen

only the *object (binary) code* which is part of the library

last compile stage *links* program with the library

Program is *independent* of underlying details

(Unix/Windows? file system type?)

*implementation* of library function can *change*

(new compiler version, bug fix, new file system)

as long as *interface* (function prototype) stays the same

## Abstract datatypes

An abstract datatype is a mathematical model for datastructures defined by the operations applicable to them (*functions*) and the constraints among them (*axioms*) without exposing details about the implementation.

ADTs *separate interface from implementation*  
the interface provides the *abstraction*  
the implementation is *encapsulated* (hidden)

ADTs allow changeable and interchangeable implementations  
client program relies only on interface, is not affected

**FILE** is an abstract datatype in the standard C library  
don't know implementation detail  
can only access with given functions (`fopen`, `fgets`, `fread`, etc.)

## Lists as abstract data types

An ADT list  $L$  with elementtype  $E$  is usually defined by:

$nil : () \rightarrow L$             empty list constructor  
can also be constant rather than function

$isempty : L \rightarrow Bool$     is empty ?

$cons : E \times L \rightarrow L$     list constructor

$head : L \rightarrow E$         head of list

$tail : L \rightarrow L$          tail of list

and the *axioms*

$$head(cons(e, l)) = e \quad \text{and} \quad tail(cons(e, l)) = l$$

## Example ADT for integer list

```
#ifndef _INTLIST_H
#define _INTLIST_H

typedef struct ilst *intlist_t;

intlist_t empty(void);
int isempty(intlist_t lst);
int head(intlist_t lst);
intlist_t tail(intlist_t lst);
intlist_t cons(int el, intlist_t tl);

// for freeing memory only: splits first element from tail
// if elp non-NULL, store value of head there
intlist_t decons(intlist_t lst, int *elp);

#endif
```

## Hiding / exposing the representation

Implementation is hidden if only a *pointer* to the data is exposed:  
incomplete structure type: `typedef struct` `ilst` `*intlist_t`  
or even a `void *` (only implementation knows what it points to)

Declaration of structure should be hidden in `.c` file  
not exposed in `.h` file (which is included by all clients)

```
struct ilst {  
    intlist_t nxt;  
    int el;  
};
```

If library client has this structure, can use internal representation  
(no longer an ADT)

## Implementing the list ADT

```
#include <stdlib.h> // for NULL and malloc
#include "intlist.h" // ensures .h and .c consistent

struct ilst {
    intlist_t nxt;
    int el;
};

intlist_t empty(void) { return NULL; }

int isempty(intlist_t lst) { return lst == NULL; }

int head(intlist_t lst) { return lst->el; }

intlist_t tail(intlist_t lst) { return lst->nxt; }
```

## Implementing the list ADT (cont'd)

```
intlist_t cons(int el, intlist_t tl)
{
    intlist_t p = malloc(sizeof(struct ilst));
    if (!p) return NULL; // could report some error
    p->el = el;
    p->nxt = tl;
    return p;
}
```

```
// returns tail, assigns *elp with head, deletes cell
intlist_t decons(intlist_t lst, int *elp)
{
    if (elp) *elp = lst->el;
    intlist_t tl = lst->nxt;
    free(lst); // just first cell, keeps rest
    return tl;
}
```



## Can we do lists of arbitrary types?

C does not have polymorphism or parametric types

⇒ cannot declare, e.g., list of *arbitrary type*

Could do: `typedef int elemtype;` (or even a `#define`)

and have everything else use `elemtype`

But need to *recompile* everything when changing `elemtype`  
binary code differs even for assignment/parameter passing  
due to varying element size; even more so for addition, etc.)

If instead of values we store *pointers* to values,  
we can have just one implementation (list of `void *`)  
must separately allocate memory for elements  
program logic must know element type (info not in the list)

## Example: list reversal in-place

Assume: we know declaration

```
struct ilst {  
    intlist_t nxt;  
    int el;  
};
```

Two pointers, splitting list:

one to part of list already reversed (initially NULL)

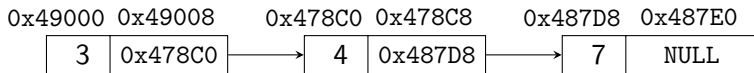
one to rest of list to be reversed (initially full list)

```
intlist_t rev2(intlist_t rest, intlist_t done) {  
    if (isempty(rest)) return done;  
    intlist_t nxt = rest->nxt; // rest to be reversed  
    rest->nxt = done; // link first cell to done part  
    return rev2(nxt, rest); // tail-recursive, becomes loop  
}  
intlist_t rev(intlist_t lst) { return rev2(lst, empty()); }
```

## Traversing linked list with address of pointer

When inserting/deleting into a linked list (e.g. *ordered* list), must change link in cell *prior* to the one inserted/deleted  
keep *address* of pointer to be changed (address of link field)  
better than with address of previous element (may not exist)

In picture, top row denotes *addresses* of individual fields



```
intlist_t hd = cons(3, cons(4, cons(7, NULL)));  
intlist_t *adr = &hd;  
adr = &(*adr)->nxt; // advance to next element
```

## Implementing a queue ADT

Queue: first-in, first-out (FIFO): insert/remove at different ends

```
#ifndef _QUEUE_H
#define _QUEUE_H

typedef struct q *queue_t;

queue_t q_new(void);
int q_isempty(queue_t q);
int q_get(queue_t q);
queue_t q_put(queue_t q, int el);
void q_del(queue_t q);
void q_print(queue_t q);

#endif
```

## Implementing a queue

To uniformly handle case when queue becomes empty/grows again we use a *dummy* cell (flag); actual first cell is *after* the dummy cell

```
typedef struct e { // cell for element, with pointer to next
    struct e *nxt;
    int el;
} elem_t;
```

```
struct q {
    elem_t *hd; // dummy; actual first cell is next
    elem_t *last; // last cell (or dummy if empty)
};
```

```
queue_t q_new(void) {
    queue_t q = malloc(sizeof(struct q));
    elem_t *p = malloc(sizeof(elem_t)); // dummy cell
    p->nxt = NULL; // no actual element
    q->hd = q->last = p; // initially both dummy cell
    return q;
}
```