Programming language design and analysis

# Domain-specific languages

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#### Part One

#### based on: Domain Specific Languages, martin-fowler.com/dslwip

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### Defining Domain Specific Languages

DSL: a computer programming language of limited expressiveness focused on a particular domain

computer programming language

used to communicate with computer *and* between humans should have fluency (beauty)

*limited expressiveness* 

can't build a software system in it contrast: data / control / abstraction structures in general PL

domain focus

makes it useful

#### Kinds of DSLs

#### external

use a different language than the application that uses them SQL, XML, awk, regular expressions (and others in UNIX)

#### internal

use same general purpose programming language as application but in a particular and limited way LISP, Ruby

#### language workbenches

IDEs for building DSLs (abstract syntax, editors, generators) more/different than usual parse/generate cycle

### Why use a DSL?

improved development productivity

communication with domain experts

change in execution context

e.g. handle definitions at runtime instead of compile time

alternative computational model not just imperative

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### What's under a DSL?

A DSL manipulates an abstraction

usually done with a *library / framework* interfaced through an *API* 

DSLs are usually a front-end to such an interface  $\Rightarrow$  the hard part is building the framework

### **DSL** Patterns

appear with internal DSLs

use syntax of underlying general purpose language for visual fluency

may need:

language with special syntactic features language where new syntax can be adapted / defined just clever use of existing syntax

## Patterns: Function Sequence

```
computer();
processor();
cores(2);
processorType(i386);
disk();
diskSize(150);
diskSize(150);
diskSize(75);
diskSpeed(7200);
diskInterface(SATA);
```

# Function Sequence: Howto

```
usually with bare function calls (global if language allows)

⇒ but needs static parsing data (context variables)

currentObject = ...
```

```
currentObject.setValue(...);
```

. . .

solution: use *object scoping* for functions and parsing data

```
computer(
  processor(
    cores(2),
    Processor.Type.i386
  ),
  disk(
    size(150)
  ),
  disk(
    size(75),
    speed(7200),
    Disk.Interface.SATA
  )
);
```

#### Nested Functions: Howto

important property: evaluation order is inside-out

(parameters before function call)

- $\Rightarrow$  good: evaluation returns fully-formed values/objects, usable further
- $\Rightarrow$  awkward: textual order is opposite to natural sequencing

Useful language features: named parameters (disk(75, 7200) is not suggestive) optional arguments variable number of arguments

### Pattern: Method Chaining

```
computer()
  .processor()
    .cores(2)
    .i386()
  .disk()
    .size(150)
  .disk()
    .size(75)
    .speed(7200)
    .sata()
  .end();
```

Modifier methods return the host object

 $\Rightarrow$  multiple modifiers can be invoked on the same object

the opposite of *command query separation* 

```
HardDrive hd = new HardDrive();
hd.setCapacity(150);
hd.setExternal(true);
hd.setSpeed(7200);
```

new HardDrive().capacity(150).external().speed(7200);

Issues:

```
naming no longer makes clear this is a setter
problems with languages where newline is a separator
finishing problem (when to stop?), esp. with nested components
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```

computer do processor do cores 2 i386 speed 2.2 end disk do size 150 end disk do size 75 speed 7200 sata end

#### end

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Express statement sub-elements of a function call by putting them into a closure in an argument.

a single Nested Closure instead of several Nested Function arguments

Issues:

needs code to evaluate the closure (vs. arguments are evaluated implicitly)

contents of closure is function sequence, still needs context variables (but they can be created before closure / destroyed afterwards) context variable can be explicit:

```
processor do |p|
p.cores 2
p.i386
end
```